Extended finite automata and decision problems for matrix semigroups

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NCMA'18

Aim: Make a connection between extended finite automata over matrix semigroups and decision problems for matrix semigroups

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• One-way finite state automaton equipped with a register

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Let M be a monoid. M-automaton is defined analogously.

• Let S be a semigroup. We want to allow the register to be multiplied elements from S.

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- We define S-automaton by letting 1 to be the identity element.

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- Let S be a semigroup. We want to allow the register to be multiplied elements from S.
- We define S-automaton by letting 1 to be the identity element.
- If S is not a monoid nor a group, then only the empty string can be accepted.

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Let S be a matrix semigroup finitely generated by a generating set of square matrices F. The **membership problem** is to decide whether or not a given matrix Y belongs to the matrix semigroup S. Let S be a matrix semigroup finitely generated by a generating set of square matrices F. The **membership problem** is to decide whether or not a given matrix Y belongs to the matrix semigroup S.

Given: $F = \{Y_1, Y_2, ..., Y_n\}$ and a matrix Y**Problem:** Determine if there exist an integer $k \ge 1$ and $i_1, i_2, ..., i_k \in \{1, ..., n\}$ such that $Y_{i_1}Y_{i_2} \cdots Y_{i_k} = Y$.

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When Y is restricted to be the identity matrix, the problem is called the **identity problem**.

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Let G be a finitely generated group and let H be a subgroup of G. **Subgroup membership problem** or generalized word problem for H in G is to decide whether or not a given element $g \in G$ belongs to the subgroup H.

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Word problem for *G* is the membership problem for the trivial group generated by 1.

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The subgroup membership problem can be seen as a special case of the (semigroup) membership problem.

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	Identity problem	Membership problem
$\mathbb{Z}^{2 \times 2}$	decidable	decidable
Н	decidable	?
$SL(3,\mathbb{Z})$?	?
$\mathbb{Z}^{3 \times 3}$?	undecidable
<i>SL</i> (4, ℤ)	undecidable	undecidable

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Theorem

Let H be a finitely generated subgroup of G. If the emptiness problem for G-automata is decidable, then the subgroup membership problem for H in G is decidable.

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- Construct *G*-automaton *V*.

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Proof.

- Let g be an element from G. We should decide whether $g \in H$.
- Construct *G*-automaton *V*.
- We are going to show that $g \in H$ iff L(V) is nonempty.

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Image: A matrix and a matrix

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Proof.

- $\{h_1, \ldots, h_n\}$ generates H
- Claim: $g \in H$ iff L(V) is nonempty.

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Proof.

• $g \in H \implies g^{-1} \in H$

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$$h_{i_1}h_{i_2}\cdots h_{i_k}=g^{-1}$$
 for some $k\geq 1$ and $i_1,i_2,\ldots,i_k\in\{1,\ldots,n\}$

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$$gg^{-1} = 1 \implies a^{k+1} \in L(V)$$

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Proof.

- Suppose L(V) is nonempty
- Acceptance condition: register is equal to identity

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$$g^{-1} \in H \implies g \in H$$

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Proof.

• Suppose that the emptiness problem for *G*-automaton is decidable.

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Proof.

- Suppose that the emptiness problem for *G*-automaton is decidable.
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Proof.

- Suppose that the emptiness problem for *G*-automaton is decidable.
- To check if $g \in H$,
 - Construct V
 - Check if L(V) is nonempty
 - \implies subgroup membership problem for H is decidable .

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Given a matrix Y from $\mathbb{Z}^{2\times 2}$ and a subgroup H of $\mathbb{Z}^{2\times 2}$, it is decidable whether Y belongs to H.

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- We will prove that the emptiness problem for Z^{2×2}-automata is decidable
- Suppose that a $\mathbb{Z}^{2\times 2}$ -automaton V is given.

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- Suppose that a $\mathbb{Z}^{2\times 2}$ -automaton V is given.
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 - ► Matrices multiplied by the register are invertible and belong to GL(2, Z)
- V is a $GL(2,\mathbb{Z})$ -automaton

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Lemma

Let G be a finitely generated group and let H be a subgroup of finite index. Any G-automaton can be converted into an H-automaton recognizing the same language.

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Lemma

Let G be a finitely generated group and let H be a subgroup of finite index. Any G-automaton can be converted into an H-automaton recognizing the same language.

Lemma

Any F_2 -automaton can be converted into a pushdown automaton recognizing the same language.

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Theorem

Given a matrix Y from $\mathbb{Z}^{2\times 2}$ and a subgroup H of $\mathbb{Z}^{2\times 2}$, it is decidable whether Y belongs to H.

Proof.

• A pushdown automaton recognizing L(V) can be constructed

- ▶ **F**₂ has finite index in *GL*(2, ℤ)
- F_2 -automaton recognizing L(V) can be constructed
- ▶ **F**₂-automaton can be converted to a pushdown automaton
- \blacksquare Emptiness problem for pda is decidable \implies Emptiness problem for $\mathbb{Z}^{2\times 2}\text{-automata}$ is decidable

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The emptiness problem and identity problem

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Proof.

• Construct an S-automaton V such that $1 \in S$ iff L(V) is nonempty.



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Undecidability of the emptiness problem for $\mathbb{Z}^{4 \times 4}$ -automata

Fact

Given a semigroup S generated by eight 4×4 integer matrices, determining whether the identity matrix belongs to S is undecidable. [KNP17]

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Undecidability of the emptiness problem for $\mathbb{Z}^{4\times 4}$ -automata

Fact

Given a semigroup S generated by eight 4×4 integer matrices, determining whether the identity matrix belongs to S is undecidable. [KNP17]

Corollary

Let S be a subsemigroup of $\mathbb{Z}^{4\times 4}$ generated by eight matrices. The emptiness problem for S-automaton is undecidable.

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Let S be a subsemigroup of $\mathbb{Z}^{4\times 4}$ generated by eight matrices. The emptiness problem for S-automaton is undecidable.

Proof.

We know that the identity problem for S is undecidable. By the above theorem the result follows.

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The universe problem and subgroup membership problem

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Let H be a finitely generated subgroup of G. If the universe problem for G-automata is decidable, then the subgroup membership problem for H in G is decidable.

Proof.

• Construct G-automaton V such that $g \in H$ iff $L(V) = \Sigma^*$.



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• Universe problem for \mathbf{F}_2 -automaton is undecidable

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- Universe problem for \mathbf{F}_{2} -automaton is undecidable
 - For a given pushdown automaton, an F₂-automaton recognizing the same language can be constructed

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- Universe problem for **F**₂-automaton is undecidable
 - For a given pushdown automaton, an F₂-automaton recognizing the same language can be constructed
 - Universe problem for pushdown automata is undecidable
- F₂ is a subgroup of SL(2, ℤ) and the membership and identity problems for SL(2, ℤ) are decidable

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We make a connection between the decidability of the subgroup membership and identity problems and the universe and emptiness problems for extended finite automata

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- Emptiness problem for S-automata
 - \blacktriangleright Decidability of the subgroup membership problem for $\mathbb{Z}^{2\times 2}$
 - ▶ Undecidability of the emptiness problem for $\mathbb{Z}^{4 \times 4}$ -automata
- Universe problem for *S*-automata

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Emptiness problem for S-automata

- \blacktriangleright Decidability of the subgroup membership problem for $\mathbb{Z}^{2\times 2}$
- ▶ Undecidability of the emptiness problem for $\mathbb{Z}^{4 \times 4}$ -automata
- Universe problem for *S*-automata

Identity and membership problems for 3×3 integer matrix groups are open.

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Thank You!

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