

# **Skew order: A new look on orthomodular lattices**

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# Orthomodular lattices I

OMLs = “Boolean algebras which are not distributive nor uniquely complemented”:

$$x \wedge (y \vee z) \neq (x \wedge y) \vee (x \wedge z),$$

$$x \vee y = 1, x \wedge y = 0 \not\Rightarrow y = x'.$$

# Orthomodular lattices II

Kröger, Beran (70's): OMLs are *skew Boolean algebras*.

$$x \dot{\wedge} y := x \wedge (x' \vee y)$$

(In Boolean algebras  $\dot{\wedge} = \wedge$ .)

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OML = “BA which are not commutative nor associative”:

$$x \dot{\wedge} y \neq y \dot{\wedge} x,$$

$$(x \dot{\wedge} y) \dot{\wedge} z \neq x \dot{\wedge} (y \dot{\wedge} z),$$

$$x \dot{\wedge} (y \vee z) = (x \dot{\wedge} y) \vee (x \dot{\wedge} z),$$

$$x \dot{\vee} y = 1, x \dot{\wedge} y = 0 \Rightarrow y = x'.$$

# Assymetry of the skew operations

Sasaki adjunction:

$$x \dot{\wedge} y \leq z \Leftrightarrow y \leq x \dot{\rightarrow} z (= x' \dot{\vee} z),$$

$$x \dot{\wedge} y \leq x,$$

$$x \leq y \Leftrightarrow x = y \dot{\wedge} x,$$

$$x \dot{\wedge} y = y \dot{\wedge} x \Leftrightarrow x C y \text{ (then } = x \wedge y).$$

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In  $\mathcal{P}(H)$  (closed subspaces/projections on a Hilbert space),  $x \dot{\leq} y$  means that

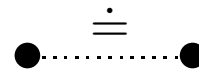
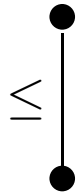
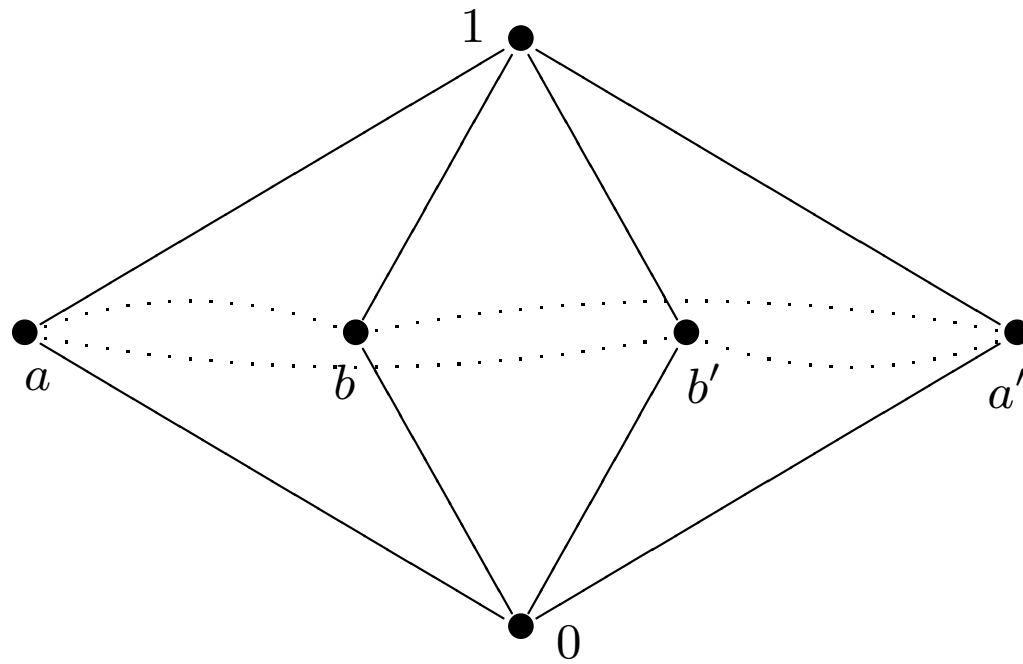
- $x$  is projected injectively into  $y$ ,
- $y$  is projected surjectively onto  $x$ ,
- $x$  is a “subspace” of  $y$  upto a deviation  $< \pi/2$ .



# Properties of the skew order

- $x \dot{\leq} y \Leftrightarrow x \wedge y' = 0 \Leftrightarrow x' \vee y = 1 \Leftrightarrow y \dot{\vee} x = y.$
- $x \dot{\wedge} y \dot{\leq} y.$
- $\dot{\leq}$  is reflexive. It is transitive or antisymmetric only in BA.
- (Skew antisymmetry)  $x \dot{\leq} y, y \dot{\leq} x \Rightarrow x = y.$
- (Skew transitivity)  $x \dot{\leq} y \dot{\leq} z$  or  $x \dot{\leq} y \leq z \Rightarrow x \dot{\leq} z.$
- $x \leq y \Leftrightarrow x \dot{\leq} y, xC'y.$  In particular,  $\dot{\leq}$  contains  $\leq.$
- When  $x \dot{=} y :\Leftrightarrow x \dot{\leq} y, y \dot{\leq} x,$  then  $(\dot{\leq}) = (\dot{=}) \circ (\leq) = (\leq) \circ (\dot{=}).$  Namely, for  $x \dot{\leq} y$  we have  $x \dot{=} x \dot{\wedge} y \leq y$  and  $x \leq y \dot{\vee} x \dot{=} y.$

# Example



# Uniqueness

**Theorem.** Let  $x\rho y$  be given by identities in language  $(\wedge, \vee, 0, 1, ')$  and variables  $x, y$  such that  $\rho$  is  $\leq$  on a Boolean algebra. Then, in any OML,  $\rho$  is either  $\leq$  or  $\dot{\leq}$ .

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*Proof.*

$$\begin{array}{ccccc}
 \{l_i = r_i\} & \xrightarrow{(\dot{=})=(\dot{\leq})\cap(\dot{\geq})} & \{l_j \leq r_j\} & \xrightarrow{\text{Sasaki adjunction}} & \{l_k \leq 0\} \\
 & & & & \downarrow \vee_k \\
 l = x \dot{\wedge} y' \text{ or } l = x \wedge y' & \xleftarrow{\text{Pavićić, Megill}} & & & l = 0
 \end{array}$$

# Skew adjointness

**Theorem.** Considering all 6 quantum meets and all 6 quantum implications, the Sasaki adjunction

$$x \dot{\wedge} y \leq z \Leftrightarrow y \leq x \dot{\rightarrow} z$$

and the *skew adjunction*

$$x \wedge y \dot{\leq} z \Leftrightarrow y \dot{\leq} x \rightarrow z$$

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(where  $x \rightarrow z = x' \vee z$ ) are the only ones valid for all OML.

*Proof.* An easy manipulation with 6-element OML

$\{0, a, a', b, b', 1\}$  eliminates other possibilities. The skew adjunction is obvious, since both sides are equivalent to  $x \wedge y \wedge z' = 0$ .

# Skew nuclei

$$j(x \wedge y) = j(x) \wedge j(y),$$

$$j^2 = j,$$

$$x \dot{\leq} j(x).$$

**Theorem.** By ordering nuclei  $j \leq k : \Leftrightarrow (\forall x)(j(x) \dot{\leq} k(x))$  we get an isomorphism  $L \cong N(L)$  between OML  $L$  and lattice  $N(L)$  of all skew nuclei.

*Proof.* Every skew nucleus is of the form  $c_a(x) = a \dot{\vee} x$  for some  $a \in L$ .

# Applications and problems

In OML,  $x \dot{\wedge} y$  is the (unique) skew largest element which is  $\leq y$  and  $\leq x$ .



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Given a (nc.)  $C^*$ -algebra  $A$ ,  $(\text{RIdl}(A), \leq, \leq)$  is a full invariant of  $A$ . TFAE:

- (i)  $\text{RIdl}(A)$  is closed under  $\dot{\wedge}$ ,
- (ii) Embedding  $\text{RIdl}(A) \rightarrow L$  is skew adjoint.

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Skew orders in non-orthomodular lattices?